

Cores of Convex Games and Pascal's Triangle

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Northwestern University
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Universidad de Vigo

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(joint with Estela Sánchez-Rodríguez)

July 4th, 2007



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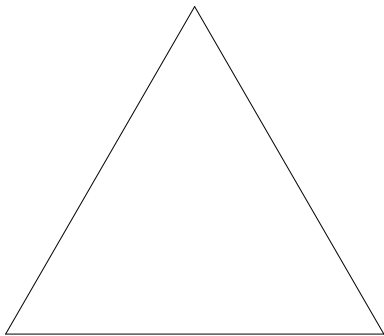
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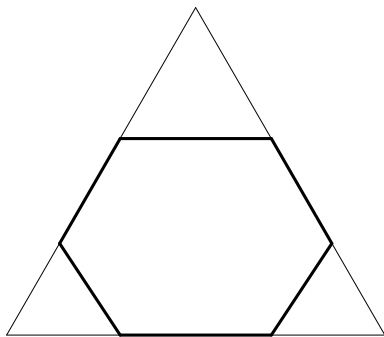
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(an $(n - 1)$ -dimensional polytope inside the set of imputations)

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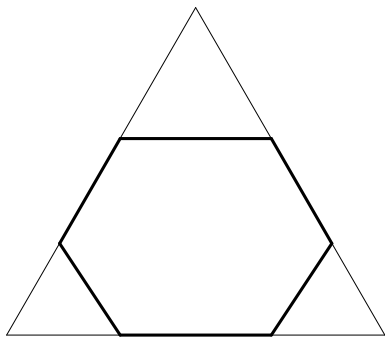


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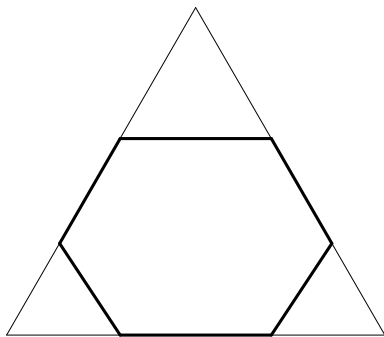
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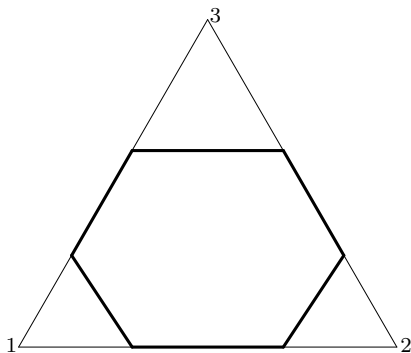
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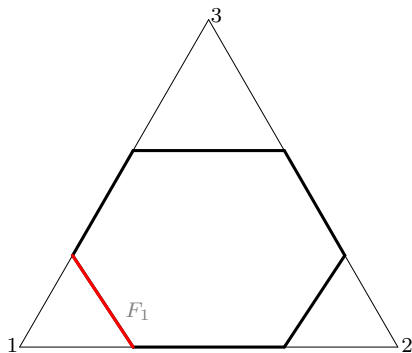
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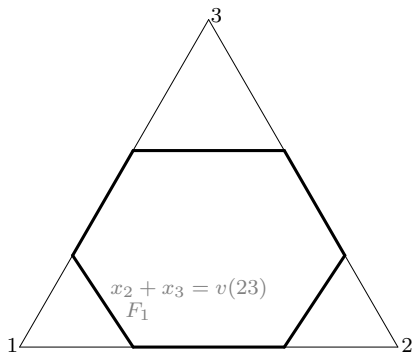
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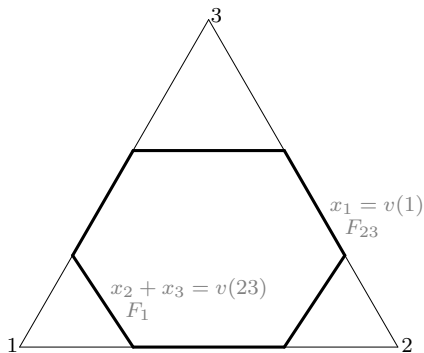
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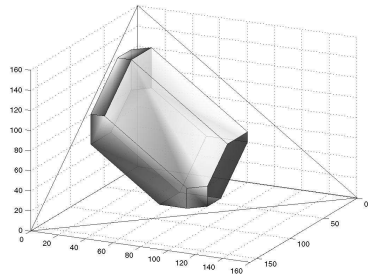
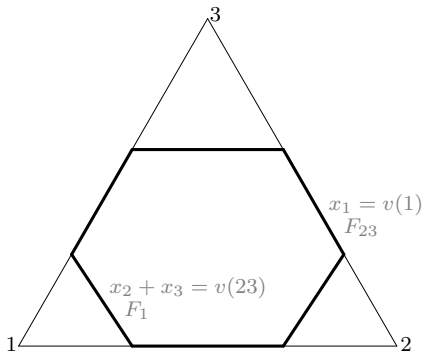
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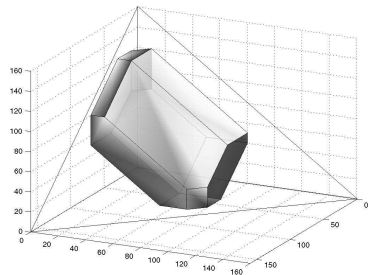
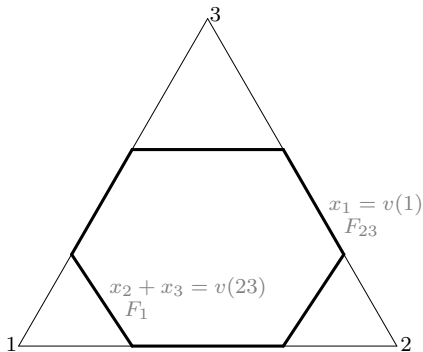
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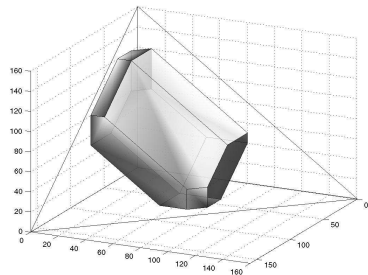
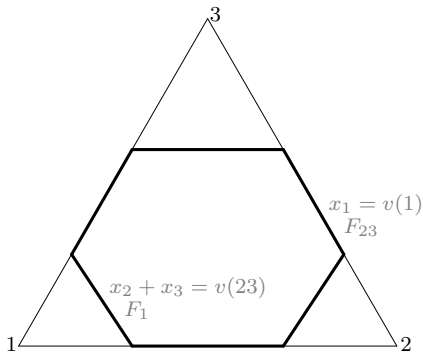


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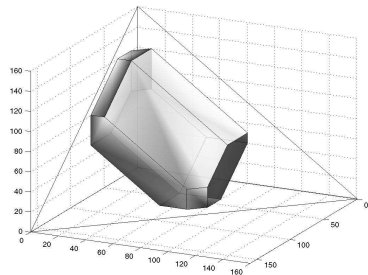
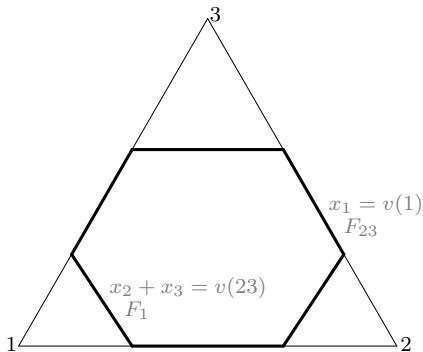
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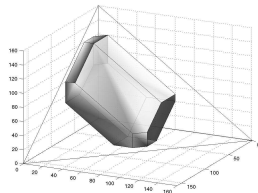
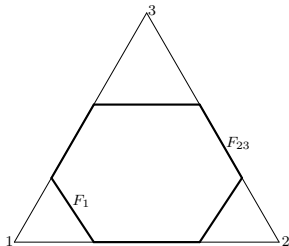
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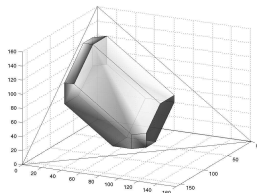
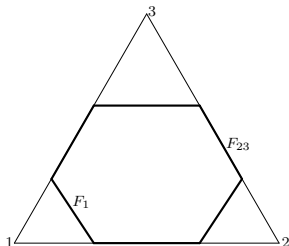
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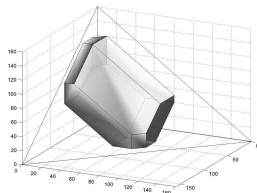
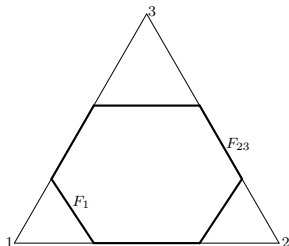


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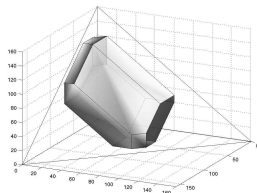
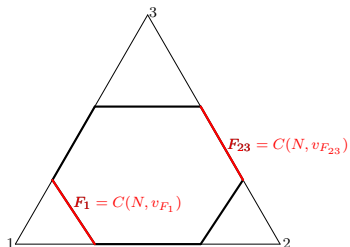


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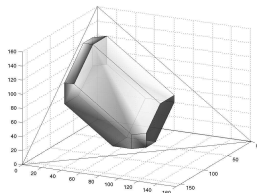
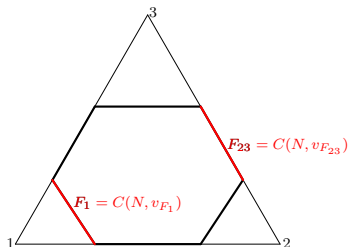


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Therefore, $C(N, v) = \text{co}\{C(N, v_{F_T}) : \emptyset \neq T \subsetneq N\}$

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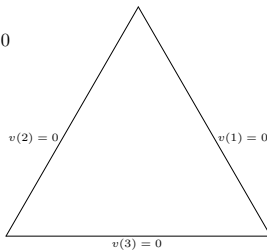
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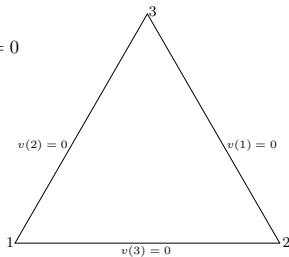
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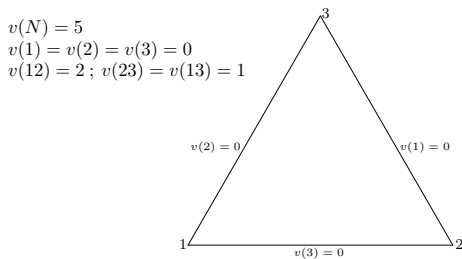


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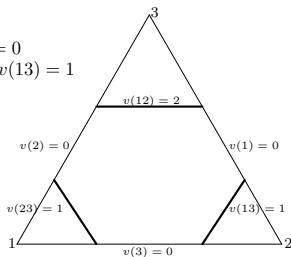


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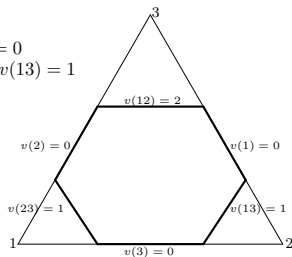
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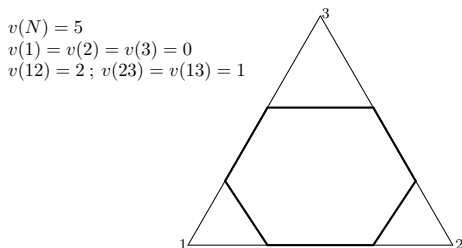


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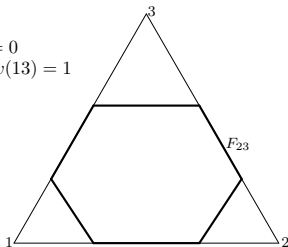


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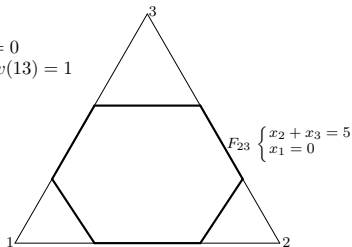


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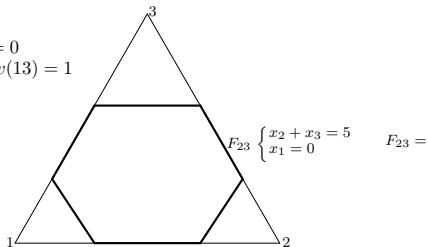


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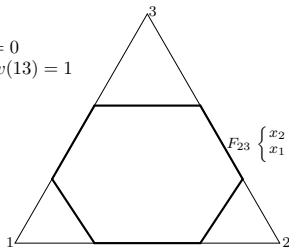


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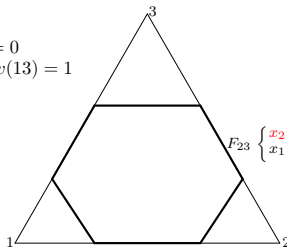
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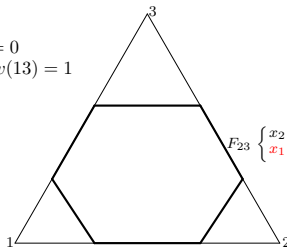
$$F_{23} = \text{red line} \times \bullet$$

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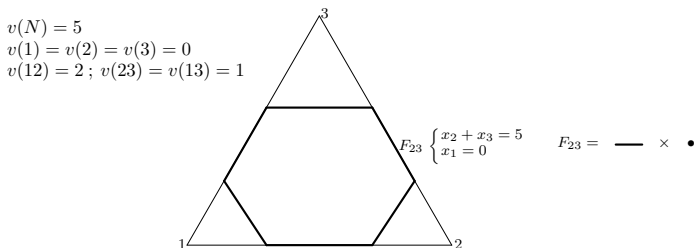


$$F_{23} \begin{cases} x_2 + x_3 = 5 \\ x_1 = 0 \end{cases}$$

$$F_{23} = \text{---} \times \bullet$$

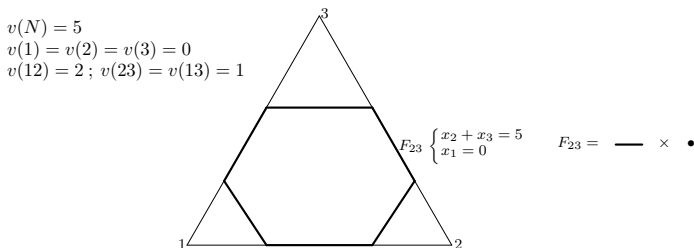
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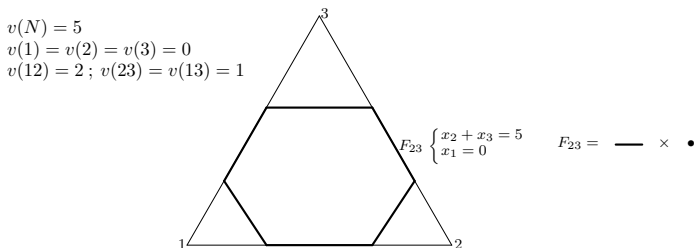
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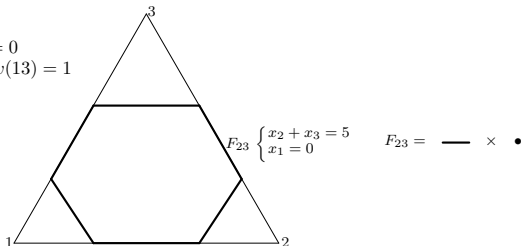
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- If $|T| > 1$, the players in T still have to “negotiate” (similarly in $(N \setminus T, v^{N \setminus T})$)

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$$v(1) = v(2) = v(3) = 0$$

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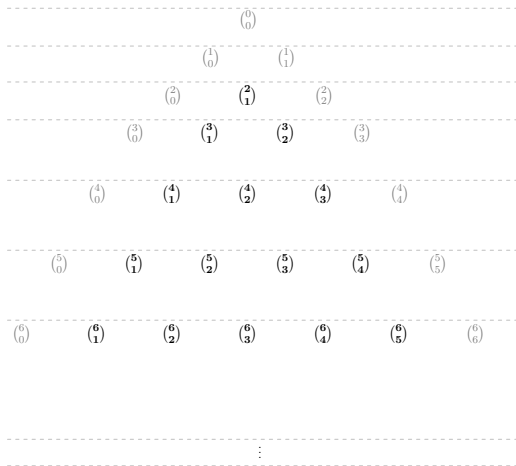
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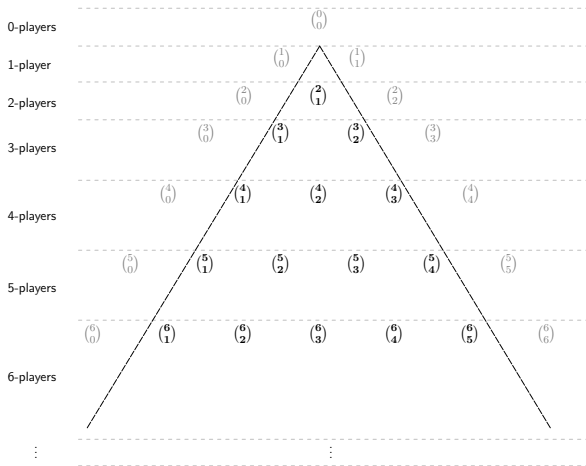
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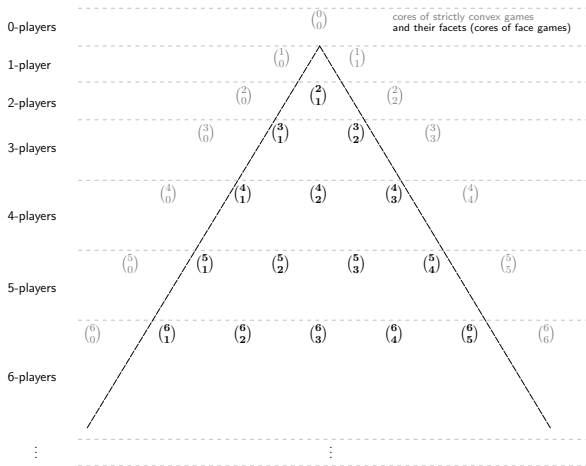
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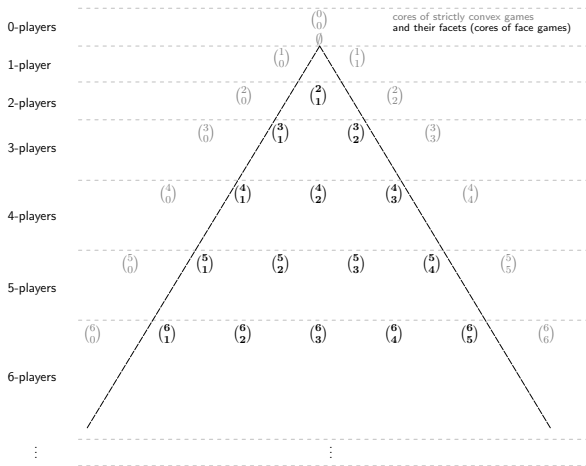
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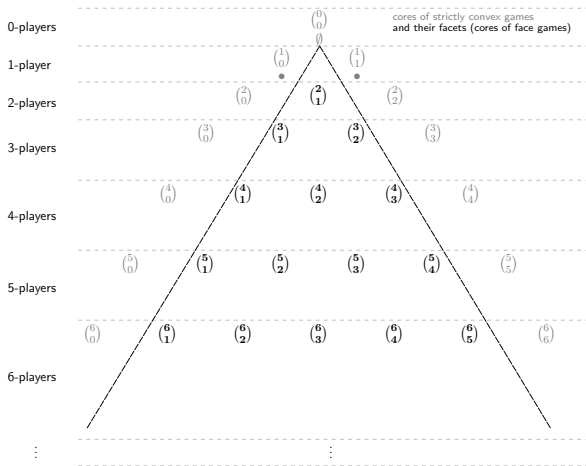
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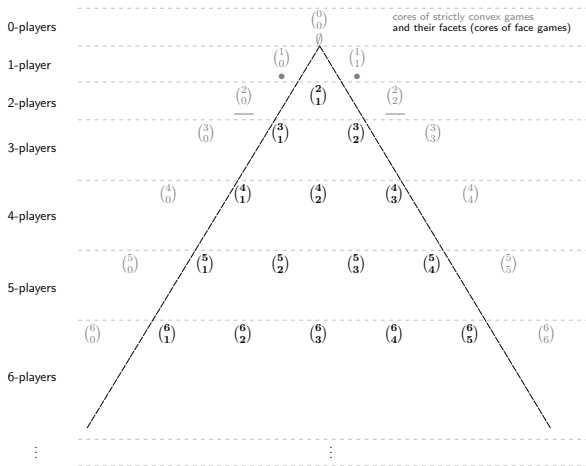
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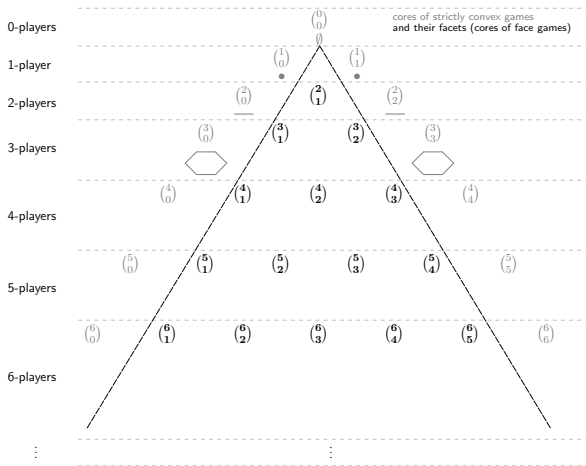
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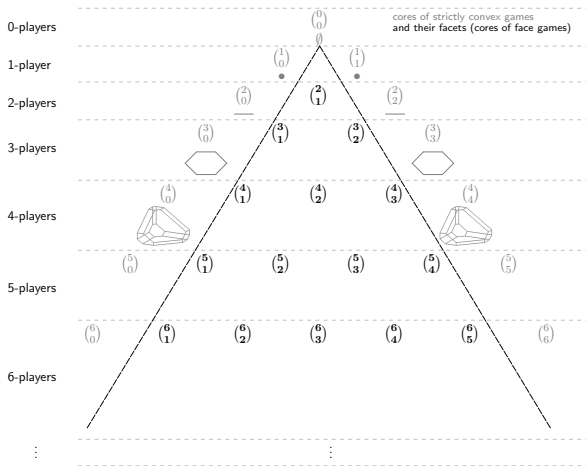
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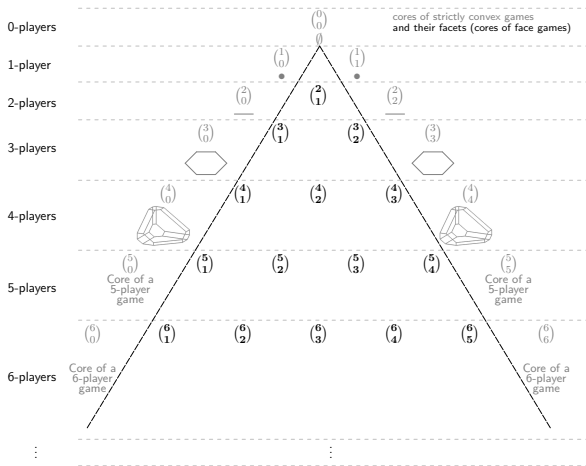
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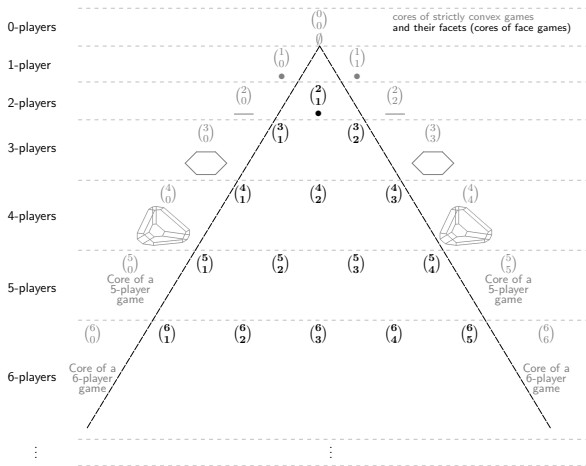
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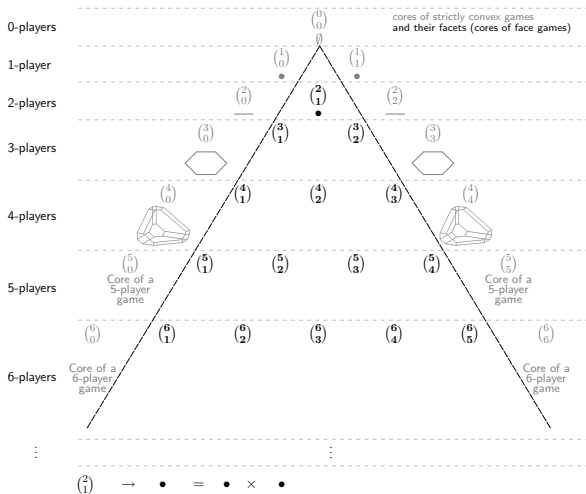
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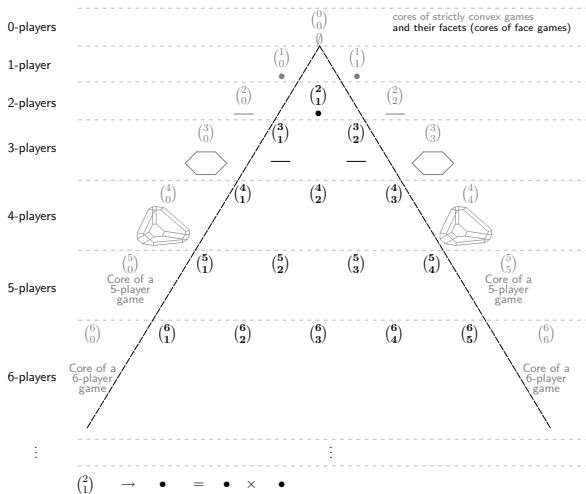
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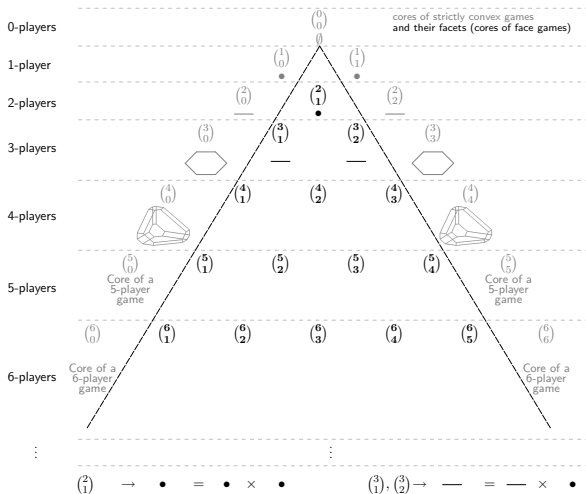
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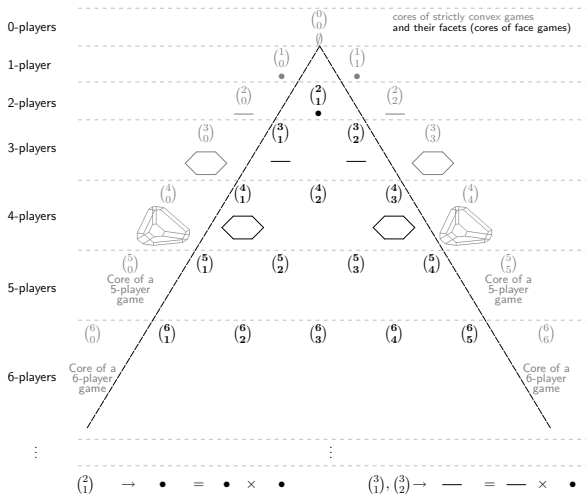
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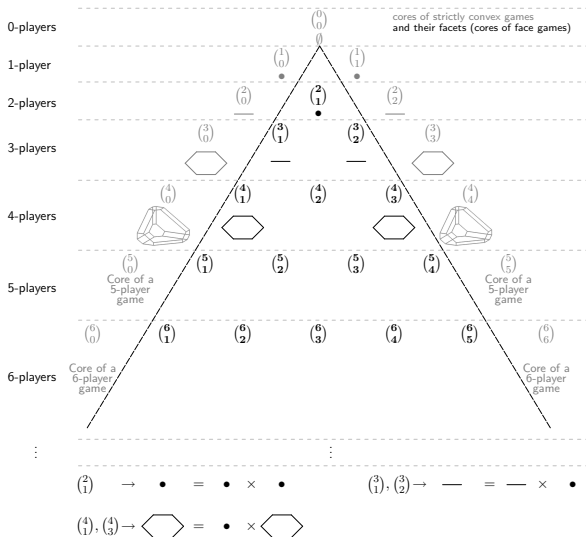
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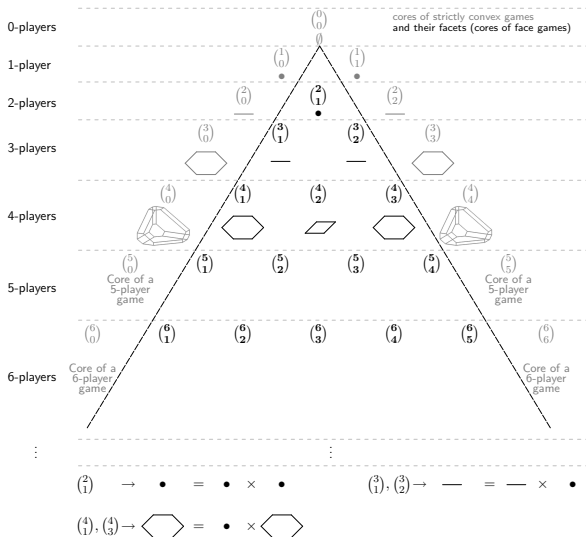
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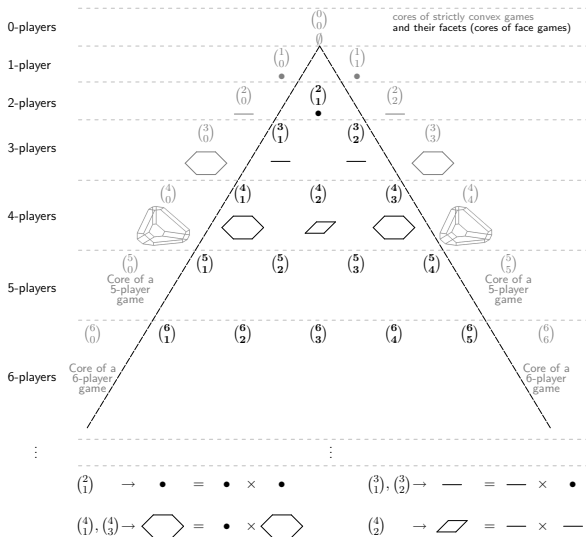
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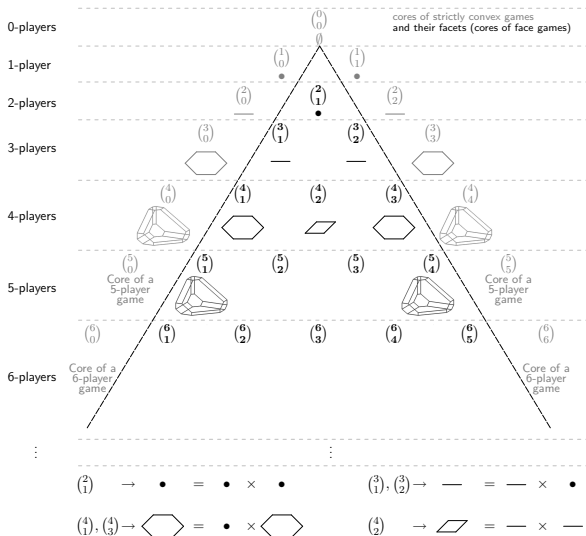
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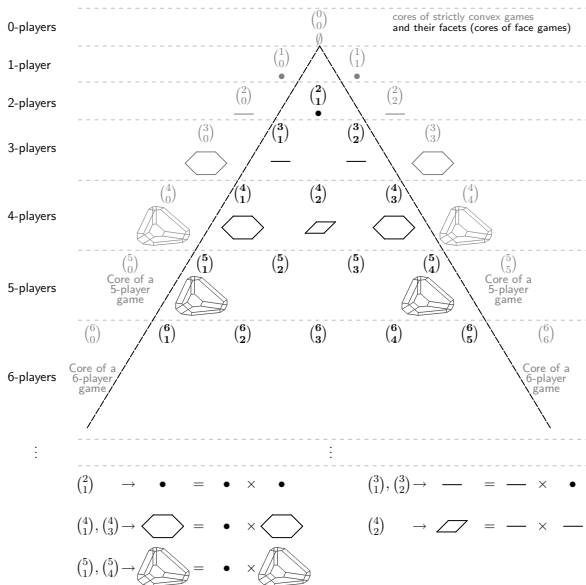
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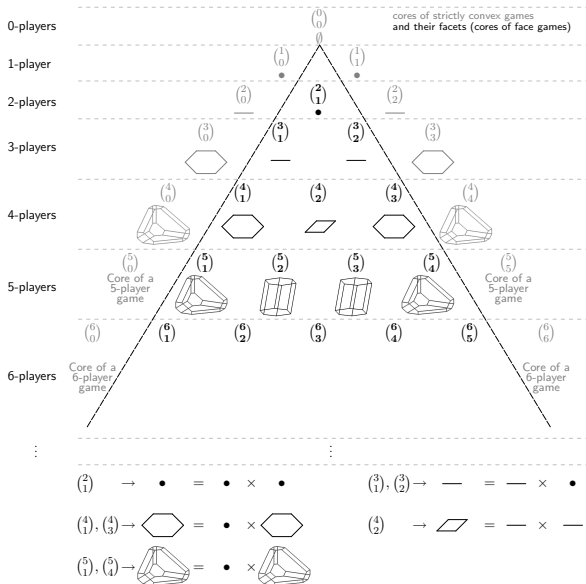
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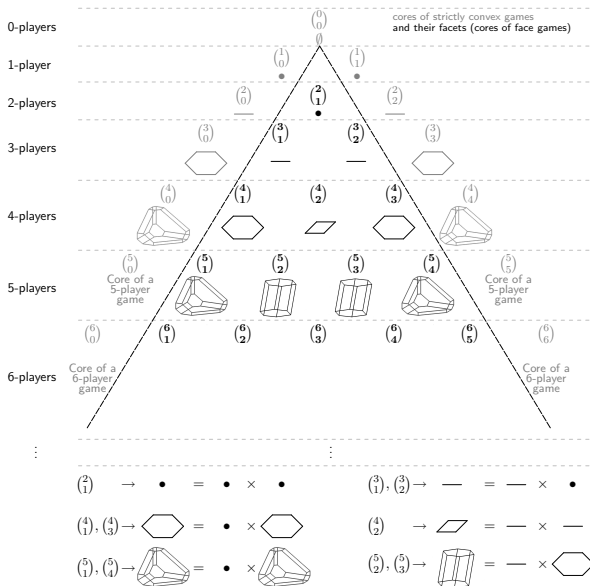
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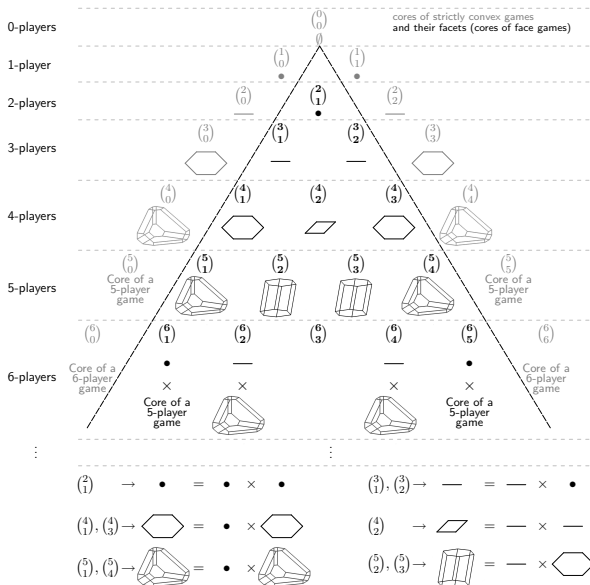
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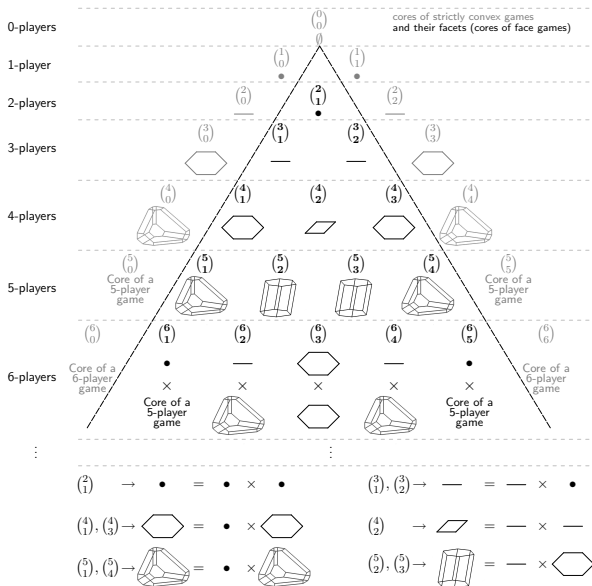
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Cores of Convex Games and Pascal's Triangle

Julio González-Díaz

Kellogg School of Management (CMS-EMS)
Northwestern University
and
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(joint with Estela Sánchez-Rodríguez)

July 4th, 2007

